

# Introduction to Algorithms and Data Structures

## 5.5. Data Structure (0): Array and Linked List

Professor Ryuhei Uehara,  
School of Information Science, JAIST, Japan.

[uehara@jaist.ac.jp](mailto:uehara@jaist.ac.jp)

<http://www.jaist.ac.jp/~uehara>

<http://www.jaist.ac.jp/~uehara/course/2020/myanmar/>

# Algorithms and Data Structure

- Algorithm: Method for solving a problem
- Data Structure:
  - Way for storing data and intermediate values
  - Influence on efficiency of a computation
  - Depending on algorithms
  - Example: array, **linked list, stack, queue, heap, tree**

# **ARRAY AND LINKED LIST**

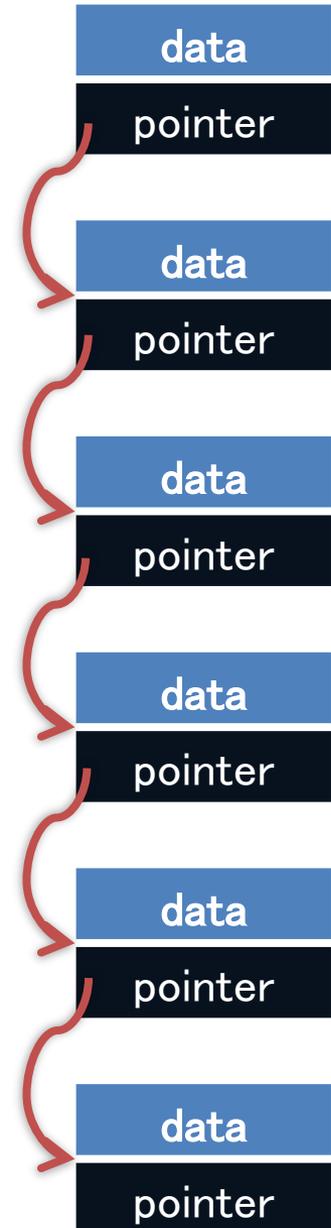
[First data structure!]

# Array: Easy to access

- Put data on **consecutive area** in memory
- We can access any index in a constant time: **Random Accessibility**
  - ↔ (Linked list does not have this property)
- Easy to access in order of indices: **Sequential accessibility**
  - ↔ (Some data may have no ordering)  
(E.g. some tree structure)

# Linked list

- Sequence of records
- They can be located in any places in memory
- Each item stores the place(s) of its “next” (and “previous”) data
  - Data: storing data
  - Pointer: indicates its “next” data (by address)
- Some variations
  - One-way linked list: Store “next”
  - Two-way linked list: Store “next”/“previous”
  - Tree can be: Store “neighbors”



# One-way linked list

- Sequence of records
  - Data: store data
  - Pointer: indicates its “next” data (by address)

```
typedef struct{
    int data;
    struct list_t *next;
} list_t;
list_t *new_r;
new_r =
(list_t *)
    malloc(sizeof(list_t));
```

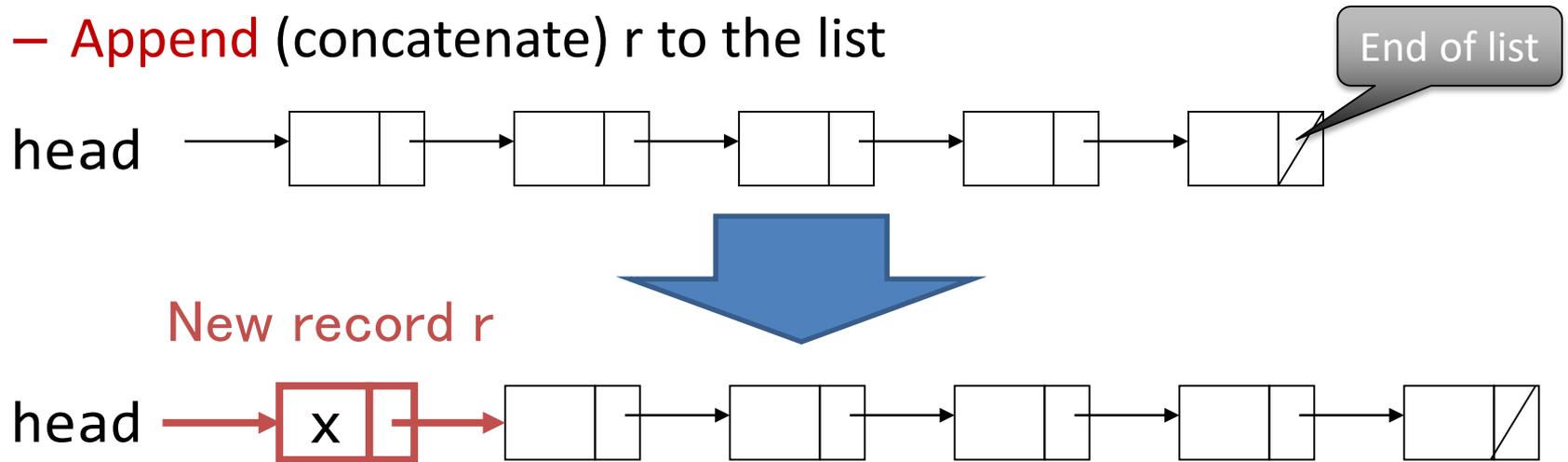
data

pointer

Initialization

# Example: Store many data into one-way linked list

- Basic:
  - Initialize list. **head** indicates **the top data**
  - For each data  $x$ , make a record  $r$  that has  $x$  in data area
  - **Append** (concatenate)  $r$  to the list

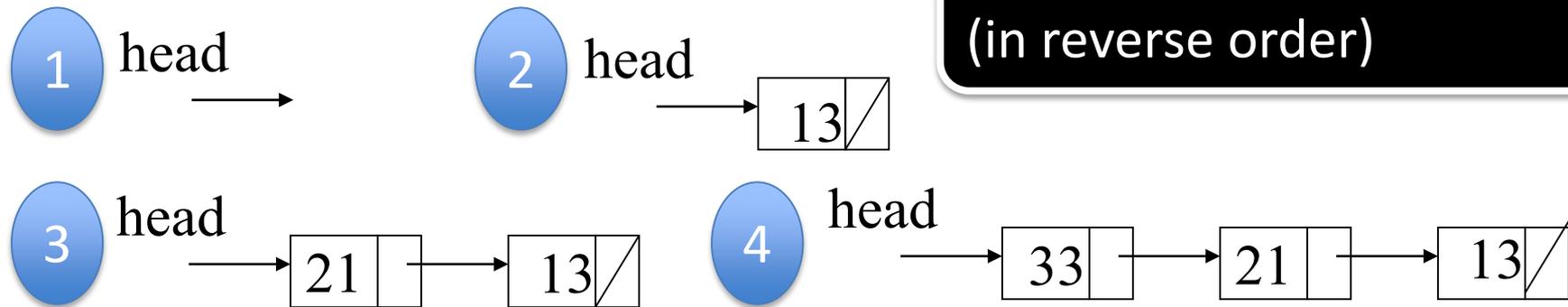


- We append at the top or last of the list

# A program for adding new records at **the top** of the linked list

```
list_t *head, *new_r;  
head = NULL;    (It is empty at first)  
while(/* we have data */){  
    int x = input data;  
    new_r = (list_t *)  
            malloc(sizeof(list_t));  
    new_r->data = x;  
    new_r->next = head;  
    head = new_r;  
}
```

New record is added at the top  
(in reverse order)

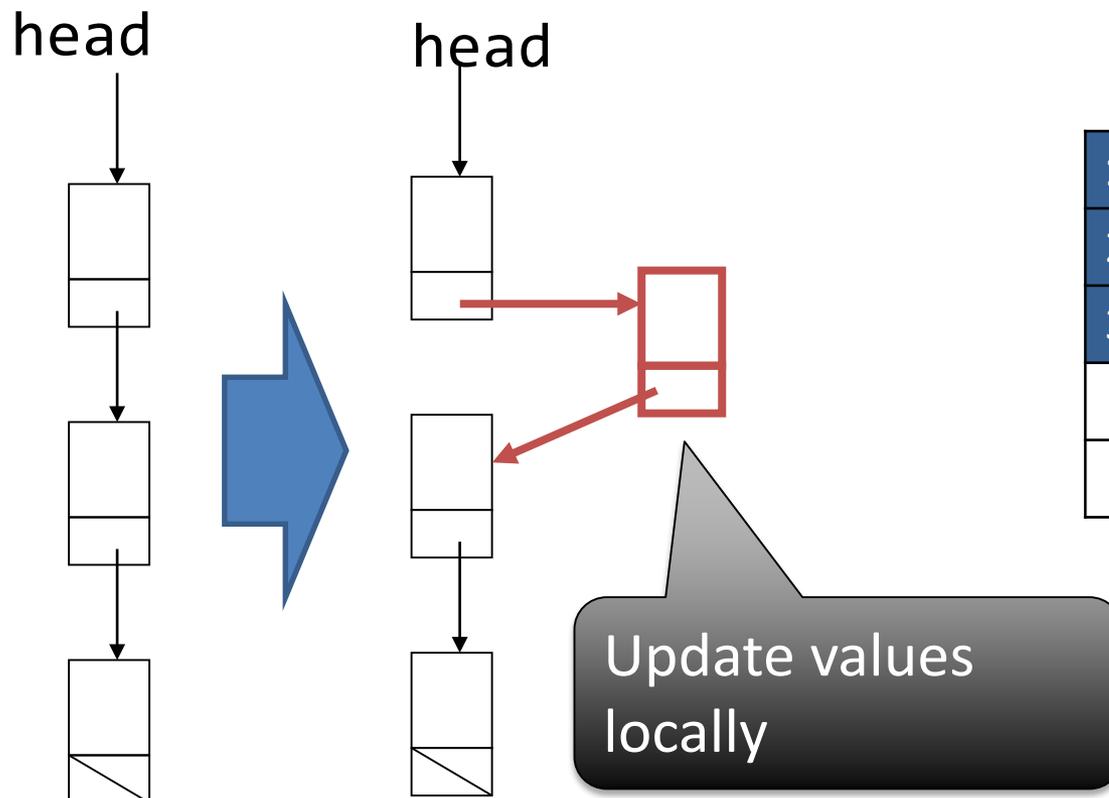


[Important point!]

# Advantage of linked list (comparing to array): Easy to add/remove data

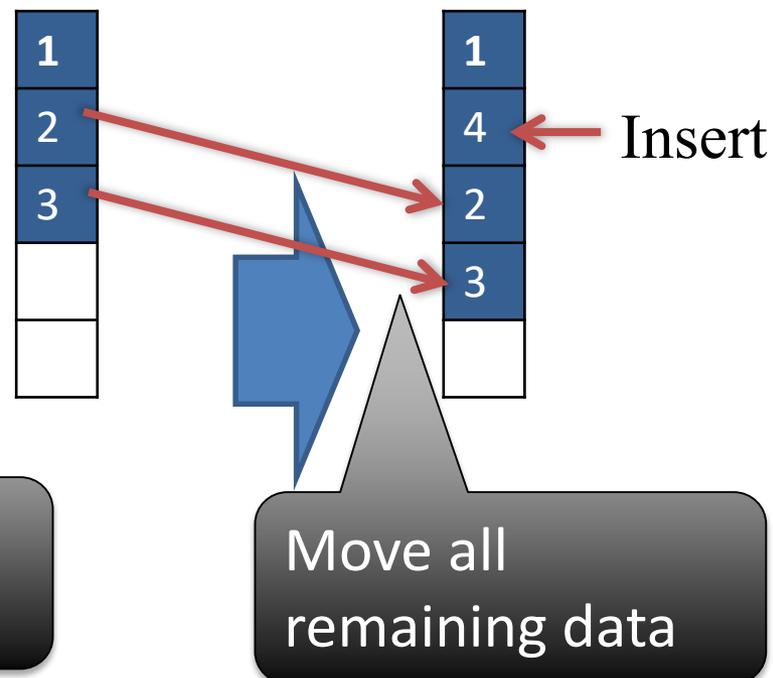
## Linked list

- No move of data



## Array

- (Many) data should be shifted



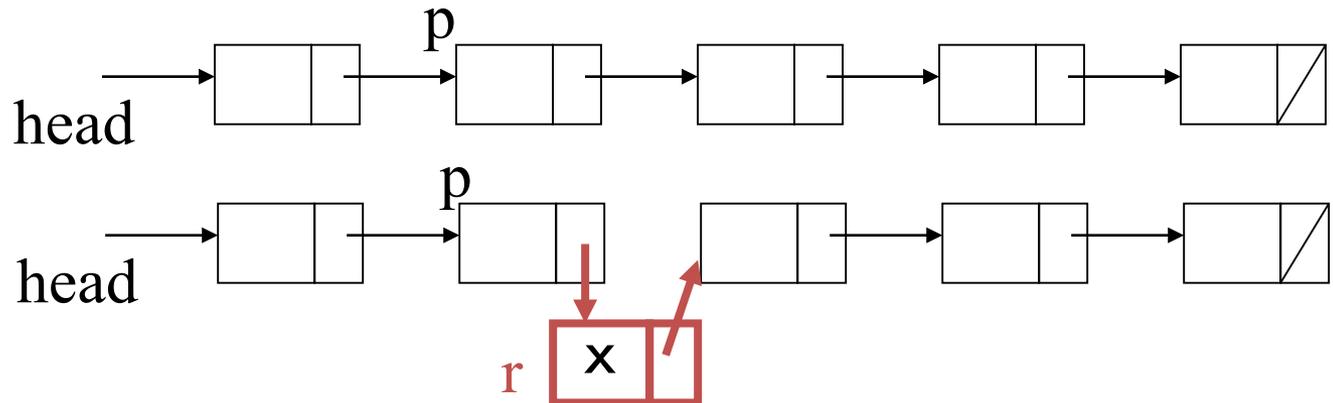
# One-way linked list:

## < Insertion of data >

- Insert data x after the node p

- Make a record r
- “Next node of r” was “next node of p”
- “next node of p” becomes r

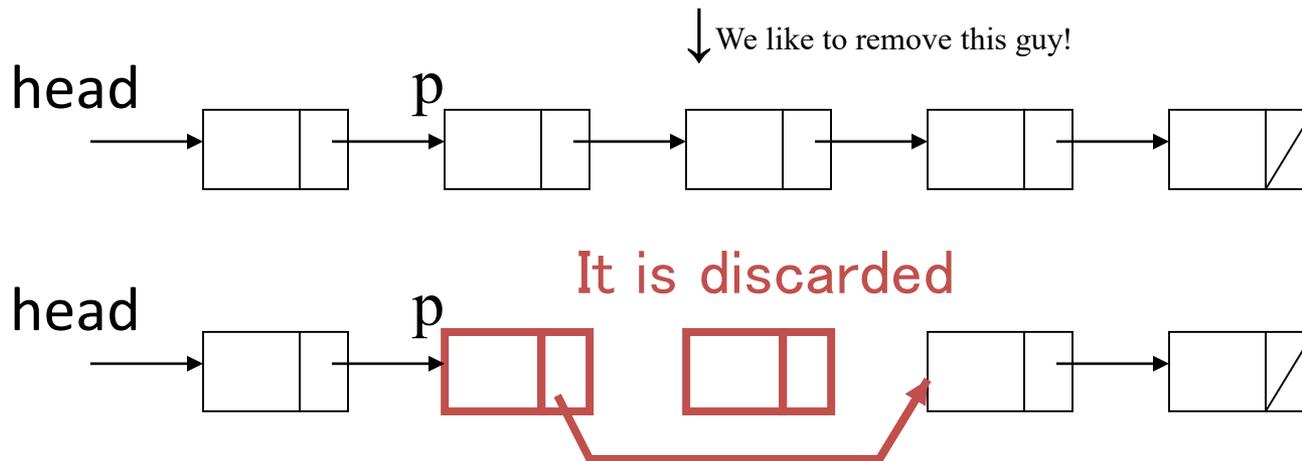
```
void insert (Node p, data x) {  
    Node r = new Node(x, p.next);  
    p.next = r;  
}
```



# One-way linked list:

## < Remove data (1) >

- Remove the **next** node of node p  
→ We can **skip** it (without removal)



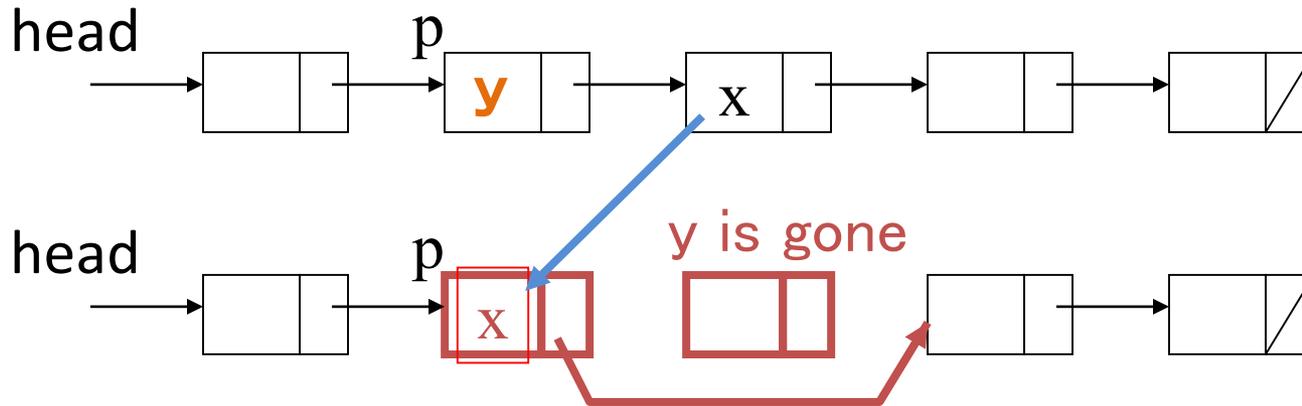
```
void deleteNextNode (Node p) {  
    p.next = p.next.next;  
}
```

[bit tricky!]

# One-way linked list:

## < Remove data (2) >

- Remove the current node p (or its **data y**)
  - We do not remove p itself, but we remove the next node q of p.  
Before removing, copy the data x from q to p.



```
void delete (Node p) {  
    p.data = p.next.data;  
    p.next = p.next.next;  
}
```

# Properties of one-way linked list

- Advantage: Easy to insert/remove  
 $O(1)$  time  $\Leftrightarrow$  array requires  $O(n)$  time
- Disadvantage: “Taking the  $i$ -th data” is slow...  
 $O(n)$  time  $\Leftrightarrow$  array requires  $O(1)$  time

Consider:

- In practical, when should you use array, and when should you use one-way linked list?
- How about advantage/disadvantage of one-way and two-way linked lists?